

## FORMULA FOR THE GENERATION OF LARGE PRIME NUMBERS

The use for large prime numbers in the 50 to 1000 digit range has become very important in recent years because of the extensive use of asymmetric (or RSA) cryptographic techniques involving public keys. The idea behind a public key  $N = PQ$  is that it is constructed by taking the product of two large primes and the knowledge that it is extremely difficult to factor such a key even when using the fastest available super-computers. At the present time the largest  $N$  which has been factored after months of effort is one of 800 bit length or about 250 digits long. This means that to keep the public key  $N$  secure one must employ primes longer than about 125 digits each. At the same time the digit length of the primes should not be more than a couple of thousand because of number crunching complications which will rise in message encoding and decoding procedures. Thus one is very interested in quickly finding prime numbers lying in the range-

$$10^{200} < P < 10^{2000}$$

How can this be accomplished? One of the oldest techniques is to generate the numbers from the Mersenne formula  $P=2^x-1$  for certain prime integers  $x$ . These numbers  $P$  are referred to as Mersenne Primes and occur only for a limited number of  $x$ s in a specified integer range for  $P$ . There are just six Mersenne primes which fall into the above stated range of interest. These are-

Exponent $x$	Digits in $P=2^x-1$
1279	386
2203	664
2281	687
3217	969
4253	1281
4423	1332

This data and other interesting facts about Mersenne Primes can be found at-

[http://en.wikipedia.org/wiki/Mersenne\\_prime](http://en.wikipedia.org/wiki/Mersenne_prime)

One of the latest found primes of the Mersenne type is  $2^{43,112,609}-1$ . The search for still larger Mersenne Primes continues but these efforts will add little to possible applications in cryptography. Since the six Mersenne Primes given in the above table are well known throughout the world, they provide little security for their full or partial use in public key generation.

As a result, one presently uses large prime numbers generated by random number generators and then tested for primality by techniques based on Fermat's Little Theorem. This works fine but can often require lengthy tests before a prime number is found. It also prevents storage and transmission of the resultant primes in the compact form of a simple formula.

We want here to discuss another approach to generating large primes based on an extension of the simple Mersenne Formula given above. The idea is to use an expansion of the form-

$$P = A + \sum_{n=1}^k c_n (b_n)^{x_n}$$

Here one chooses the values of the coefficients  $c_n$ , the base  $b_n$ , and the exponent  $x_n$  beforehand to make  $P$  fall into the desired digit length range. Once this has been accomplished one next varies the value of  $A$  until  $P$  is found to be prime. This procedure will produce a large group of primes in any desired digit range. It also allows one to store and transmit large primes  $P$  as the simple expression-

$$P = P[A, (c_1, c_2, c_3, \dots), (b_1, b_2, b_3, \dots), (x_1, x_2, x_3, \dots)]$$

For example, one can write by this method the primes  $P=13110$  and  $P=531457$  as  $P[29, (1), (2), (17)]$  and  $P[16, (1), (3), (12)]$ , respectively. Here is a generated 90 digit prime-

$$P[80, (3, -4), (5, 9), (127, 67)] = 80 + 3(5^{127}) - 4(9^{67}) =$$

1763241526233431261953104462031586229870913053609478643474083  
26824097585841856401800996179

and a 336 digit prime-

$P[-101,(1,-15),(3,13),(703,219)]=$

2607606577959774671829359348511190678090222551929747636956330  
2149423993280957825214908265755063915887047561169170339029906  
3818367294333854126397918281127482463231778303637616188140499  
6834954099445823562592683135365418762645362428872494045936481  
3711457793470404678406881576035539790976839999794852198869613  
2557773101732783986986203857973

It took me just 10 minutes to generate this last prime by carrying out the computer operation-

```
V:=3^703-15*13^219;  
for n from -150 to +150 do {n,isprime(V+n)}od;
```

It picked up a **true** response for  $n=-101$ .

To establish the size of the prime number one first evaluates the sum P-A. Lets say I want a prime number which is 200 digits long. Then I first make a random try –

$$P-A=3*6^{234}+5*9^{129}$$

This evaluates to a 188 digit long number. Hence I raise the exponent 234 to 256 and one now has a 200 digit number. Next I use the MAPLE program-

```
W:=3*6^256+5*9^129  
For n from -m to +m do {n,isprime(W+n)}od
```

This time it took a me about 15 minutes evaluating things in 200 unit chunks until the computer spit out  $n=2896$  as **true**. Hence we find, among numerous other possible 200 digit long primes, the prime-

$P[2896,(3,5),(6,9),(256,129)]=$

4828823736718668453618529174073191219493074044100855325961238  
 1139792832987949909815299583440559623645889548948958730003394  
 5951647082547518376016369772693029222320269307492769462133629  
 35340793610301309

It should be pointed out from the basic Prime Number Theorem that the value of A can become quite large when P-A is large. One knows that the number of expected primes up to a number N is approximately N/ln(N) and that up to a number M it is M/ln(M). Hence a simple calculation suggests that the distance between neighboring primes will be about -

$$\Delta s = \frac{[M - N]}{\left[ \frac{M}{\ln(M)} - \frac{N}{\ln(N)} \right]}$$

If we now take M=1.1N and N=10<sup>m</sup>, then we find-

$$\Delta s \approx m \ln(10) \text{ for } m \gg 1$$

One should expect the first value of A to be comparable in magnitude to Δs. Thus for the 175 digit long number-

$$P = A + 5^{237} + 11^{73} + 3 \cdot 13^{156}$$

We find in the range -7000<A<7000 that P will be prime if A=-6312, -4890, -3806, -1302, 172, 2652, 2728, 4770, and 6634. Thus the spacing between neighboring As is about 13,000/9≈1440. The above Δs approximation gives a somewhat lower estimate of Δs=175 ln(10)=403. The Δs criterion can therefore be thought as only an estimate for the number of integers in the separation of A values lying between prime values of P.

The number of large primes which can be generated by the present method are very large (actually infinite) and should be able to play an important role in the quick production of public keys for RSA cryptography. This is because an adversary would not be able to come up with the multiple combinations of A, c, b, and x which can be used to generate a given prime

P. It might even be safe to transmit a coded version of one of these larger  $P_S$  to a friendly receiver and he in turn could send a return coded version of his own generated  $P_R$ . This would place into the hands of both the sender and receiver the ability to generate a public key  $N=P_R P_S$  which will not be breakable by a third party. Secure encoded messages could then be sent either way by the sender and receiver without anyone else being able to decipher things.

Here is a final example of a large prime number generated by the above method. First we pick at random the number-

$$P-A=(3*7^{317})+(3*11^{409})+(13^{283})$$

Here we are dealing with a 427 digit number. A simple evaluation for different values of A near zero leads to the value  $A=-240$  needed to make P prime. Other values found in the  $-3000 < A < 20,000$  range are  $A=-2910, 1786, 4758, 7632, 10558, 13782, 16626, \text{ and } 19392$ . We were aided in finding these other values of A by noting that  $\Delta s=427 \ln(10)=983$ . The actual distance between neighboring As is about 2800. Using  $A=-240$ , we have the prime-

$$P[-240,(3,3,1),(7,11,13),(317,409,283)]=$$

2551111758131785618119558085058829315239823557156036734365596  
 1546638739124033630305658119899856751546626372320279086027937  
 0842643010137646603230957468474772113623714050629234533966093  
 4090715631035805666578978411383443196665793318493125853190517  
 7956447776135896231377773976004118570590211271016814150751329  
 0109962237164039569152249163286095425765269456086220714687249  
 6194934490405048750131991699976880775333182190993194511309051

Taking the product of this last prime with the smaller 176 digit prime-

$$P[-576,(1,1,3),(5,11,13),(239,73,157)]=$$

2323954222420997562680523557530151904666756345167574015067027  
9416746400377132053770925168581502203863219678986785875947586  
706231777043739499449465932122627413989424782093519279

leads to the Public Key-

PubKey =

5928666942178217851926933781782541665051601566122142190040100  
4496681809372387577488824679881737821843798709239326900641649  
7817974198482782480974479048374790013821022133919589702819575  
4147737209941444602703724279511854121946500001607107061673232  
2351063110978296410203728235639521596469054186439058250559146  
1931338593341826342198954307919878451233532347219662591438289  
7957138292003063502660452452634253397527194895813243884310122  
9481507971530401419000766322540797331494848456559738956367532  
8808651296756820762665436217372639654236137304426496402763676  
53584564931517270464746004193727941164906261795694229

If a code breaker were not aware of either one of the two Ps used in this last product, there would be no way at present he could successfully factor this 602 digit Public Key.

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